DEPARTMENT OF COMPUTER SCIENCE



CSCI-564 CONSTRAINT PROCESSING AND HEURISTIC SEARCH

LECTURE 16 – ADVERSARY SEARCH

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We can represent most of problems as a shortest path problem.





Heavy computation.



g

goal



We can represent most of problems as a shortest path problem.



Help the search effort, but some problems can be too large.



Abstraction of the problem UNIVERSITY AS IT'S MEANT TO BE







Or we can use **Real-Time search** algorithm.

- Suboptimal solution
- Decide the time allowed to the search between each action.





- We didn't discuss problem with adversaries.
- How can we model problem in which you compete with another agent?
 - One way is to represent the problem as a game.
 - Usually referred as Game Theory.

- What difference does it make to search of the optimal solution?
 - Adversaries introduce uncertainty.
- You don't decide of all the actions in the search tree.



- In game theory, we are not talking of optimal solution.
- We want to find the optimal strategy.
- What is the difference?
 - Optimal strategies result in perfect play.
 - The players take actions alternately and independently.



- We will focus on algorithms such as negmax and minimax.
- With one pruning strategy: $\alpha\beta$.
- In game theory:
 - The search trees are rather depth-bounded than cost-bounded.
 - The value are computed with a static evaluation function.
 - Why?

Nondeterministic environments

- Not only problems with adversary agents can be represented this way.
- In nondeterministic or probabilistic environments:
 - We include problem where the "adversary" is the unpredictable behavior of nature.
 - The outcome of an executed action in a state is not unique.
 - The ack of knowledge for modeling the real world precisely.
 - Sensors and actuators that are imprecise.
 - Etc.



Nondeterministic environments

- Solutions to nondeterministic problems are not sequences of actions.
 - Why?
- Solutions are presented as mappings from state to actions.
 - We call it policies.
- It requires state space traversal to return a solution.
 - Policy are often represented as value function.
 - Assigns a value to each state.

Two-Player Games

StFX

- Chess is the typical example for two-player games.
- One of the main successes in AI.





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Two-Player Games

- To select an optimal move in a two-player game, we construct a game tree.
 - A node represent a board configuration.
 - The root is the current configuration.
 - The children are reachable configurations.



Two-Player Games

- Game theory is an entire field of research.
- We will focus on zero-sum games.
 - The win of one player is the loss of the other.
- And games with perfect information.

- The tree is generally too large in two-player games.
- Terminal nodes are evaluated by a heuristic procedure called static evaluation function.
- Making decision before knowing the optimal action is similar to real-time search.





• Static evaluator:

- Doesn't need to be correct.
- But it needs to yield the correct values for terminal positions (Goals).
- And that higher values correlate with better positions.
- There is no notion of admissibility.



STFX

- The performance of the algorithm depends on the static evaluation function.
- What is a good static evaluation function?
 - Balance between giving an indicative value without excessive computational cost.
 - Usually developed by expert in laborious, meticulous trial-and-error experimentation.
- For complex problem, a perfect evaluation function doesn't exist.

- A first algorithm: negmax.
- The idea is simple:

- We assume that the value of a position for the first agent is the negative of its value for the second agent.
- Meaning that we choose a move that maximizes his worst-case return.

StFX

Procedure Negmax Input: Position *u* **Output:** Value at root

if (leaf(u)) return Eval(u) $res \leftarrow -\infty$ for each $v \in Succ(u)$ $res \leftarrow max\{res, -Negmax(v)\}$ return res ;; No successor, static evaluation ;; Initialize value *res* for current frame ;; Traverse successor list ;; Update value *res* ;; Return final evaluation





Procedure Negmax Input: Position *u* **Output:** Value at root

StFX



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Procedure Negmax Input: Position *u* **Output:** Value at root

StFX

if (*leaf*(*u*)) **return** *Eval*(*u*) ;; No successor, static evaluation ;; Initialize value res for current frame $res \leftarrow -\infty$ 6 ;; Traverse successor list for each $v \in Succ(u)$;; Update value res $res \leftarrow max\{res, -Negmax(v)\}$;; Return final evaluation return res -4 -6 What do you think about this result? 6 9 4 -5 -3 -9 -6 0 -2

- A second algorithm: minimax
- The idea is similar:
 - The tree consists of two different types of nodes.
 - The MIN nodes for the player that tries to minimize the payoff.
 - The MAX nodes for the players that tries to maximize the payoff.
 - The agent try to maximize its payoff during its turn.
 - And the adversary try to minimize the payoff of the agent during his turn.

StFX

Procedure Minimax Input: Position *u* **Output:** Value at root

```
if (leaf(u)) return Eval(u)

if (max-node(u)) val \leftarrow -\infty

else val \leftarrow +\infty

for each v \in Succ(u)

if (max-node(u)) val \leftarrow max\{val, Minimax(v)\}

else val \leftarrow min\{val, Minimax(v)\}

return res
```

;; No successor, static evaluation ;; Initialize return value for MAX node ;; Initialize return value for MIN node ;; Traverse successor list ;; Recursive call at MAX node ;; Recursive call at MIN node ;; Return final evaluation

Procedure Minimax Input: Position *u* **Output:** Value at root

StFX

if (leaf(u)) return Eval(u)if (max-node(u)) $val \leftarrow -\infty$ else $val \leftarrow +\infty$ for each $v \in Succ(u)$ if (max-node(u)) $val \leftarrow max\{val, Minimax(v)\}$ else $val \leftarrow min\{val, Minimax(v)\}$ return res

;; No successor, static evaluation ;; Initialize return value for MAX node ;; Initialize return value for MIN node ;; Traverse successor list ;; Recursive call at MAX node MAX nodes ;; Recursive call at MIN node ;; Return final evaluation MIN nodes MAX nodes **MIN** nodes 6 5 3 8 9 0

Procedure Minimax Input: Position *u* **Output:** Value at root

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What do you think about this result?

Game Search Tree

- The game tree cannot be fully evaluated.
 - So, we need to define a maximum depth.
- In practice, the depth that can be explored depends on a time limit.
 - The computation time is not known beforehand.
- Usually, you apply an iterative-deepening approach.
 - You increase the depth by 2 until the time available is exhausted.